

## **AUTOMATION SYSTEM AND METHOD FOR OPERATING AN AUTOMATION SYSTEM**

This is a Continuation of International Application PCT/EP99/06902, with an international filing date of September 15, 1999, the disclosure of which is incorporated into this application by reference.

### **FIELD OF AND BACKGROUND OF THE INVENTION**

5           To start-up and service an automation device, monitoring of the software running on the automation device is provided. It is desirable that said monitoring does not noticeably interfere with, or even interrupt execution of the software.

### **OBJECTS OF THE INVENTION**

10           Thus, an object of the present invention is to provide an automation system, and a method for operating the automation system, in which an automation device can be monitored by a programming device such that interference is avoided to the greatest possible extent.

### **SUMMARY OF THE INVENTION**

15           This and other objects of the invention are attained by an automation system, and method for operating the automation system, that comprises at least one automation device and one programming device connected to the automation device, enabling the two devices to communicate with one another. The automation device executes an application program, in  
20           particular, a program to control and/or monitor an external technical process.



device, where they are executed under the control of the firmware of the automation device. Each automation device and programming device can include a computer readable medium, on which is stored a computer program for operating and monitoring the automation system.

5           In automation tasks, processing is typically cyclical. For example, a few modules that are continuously restarted in a fixed time pattern. Some special modules are also restarted immediately after execution. This is termed a free cycle, i.e., a cycle that is not based on a fixed time pattern.

10           If the automation device has several processors, several modules can be executed simultaneously. If only one processor is available, however, the corresponding modules are executed in a time-sharing operation, which gives the user the impression that the corresponding modules are processed simultaneously.

15           The modules are created by the user in languages provided for this purpose (possibly several different languages) and are translated into the machine language of the automation device by a translation program executable on the programming device. In an exemplary embodiment, the programming languages KOP, FUP, AWL, SCL, HIGRAPH, GRAPH 7, etc. are used.

20           The automation device successively executes the individual machine instructions of a module. To this end, each individual machine instruction of the automation device can be uniquely referenced by a code address. This

code address is specified by the module, the module type and number, and the individual machine instruction, which is represented in the module by its corresponding address. The aforementioned request message comprises one or several of said code addresses.

5           To select a data address or a data address area, a certain instruction of the application program or a section with several program instructions is flagged and then the flag is interpreted as a selection. Alternatively, this selection can of course also be made by explicit data, for example by manually entering the relevant data addresses.

10           In an automation device, data that is supplied by, for example, sensors is processed and actuators are initiated based on the processing results. There are different constructs for this purpose depending on the language used. They all have in common that they are capable of differentiating between various data areas, which they address in their instructions. In all languages, it  
15 is possible to poll an individual sensor, address an individual actuator, and temporarily store the results, e.g., to retain states across module boundaries. Depending on the machine proximity of the correspondingly used language, the language scope also supports a direct manipulation of registers. The smallest unit thus addressable is a data address.

20           To monitor the automation device, the programming device sends a job to the automation device by means of a request message. This request

message comprises one or several code addresses corresponding to one or several data addresses per code address.

The request message has a header segment, an entry data set segment, and one or several watch point segments, each of said watch point segments  
5 having at least one code address and one or several data addresses. The data addresses can address all memory areas existing in the machine model, including the registers, but particularly bit groups, since many sensors and actuators supply not only individual bits but also groups of bits. For this reason, larger connected areas may also be addressed with a data address.

10 When reaching one of the code addresses specified in the request message, the automation device records the data indicated in the same watch point segment together with this code address. In addition, a trigger condition may be defined, which must be met when reaching a code address, so that recording will start. The automation device in turn transmits the requested  
15 data to the programming device after each cycle, regulated by flow control, until the job is stopped or cancelled by other request messages.

The structure and form of the result message are entirely analogous to that of the request message. The result message also has a header segment, an entry data set segment, and one or several watch point segments, each of said  
20 watch point segments having at least one code address and one or several data addresses.

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The constructs available in the so-called higher programming languages are mapped to the instructions of the target machine. In some situations, certain data may be available in the corresponding registers of the automation device only temporarily without ever being stored under a data  
5 address that can be referenced by the user. To enable monitoring of such interim results on the side of the programming devices, the request message of the present invention can indicate registers to a code address. As an example is a simple logic function in the form of "IF sensor 1 AND sensor 2 OR sensor 3, THEN actuator 1." The result of the AND operation is  
10 intermediately stored in a register so that this intermediate result can be used in the subsequent OR operation. The ultimate result of the logic operation is then again directly available at a data address as the content of the variable actuator 1. The intermediate result, however, would be lost if access to the register itself were not also possible. At this point it is also readily evident  
15 why code addresses are utilized in the present invention in the request message as well as data addresses for each of these code addresses.

Since it is perfectly possible in modules that a data address is written and read several times, it is advantageous to monitor the data not at the end of the cycle, but at all the code addresses at which it is referenced. This makes it  
20 possible to reconstruct the application program on the side of the programming device.

Particularly in the case described above, it is not sufficient, for instance, to indicate the register in which the intermediate result of the aforementioned AND operation is buffered, since this register may be used to store a wide variety of intermediate results at various instants. Only  
5 immediately after the logic function "IF sensor 1 AND sensor 2," the content of the register reflects the result of the AND operation. That is, the result of the subsequent OR operation is buffered in the same register, so that the result of the AND operation is lost.

In operation, the automation device records the data specified in the  
10 request message in order to send it to a programming device subsequent to recording. Generally, the automation device can record only a comparatively small amount of data. In particular, the amount of data to be recorded is defined on the one hand by the type and speed of the automation device, and on the other hand by the communication means between the automation  
15 device and the programming device as well as by the interference tolerated by the user in program execution on the side of the automation device. The amount of transferable data defines an upper limit for the amount of recordable data, even if the automation device could theoretically record a larger amount of data. It is therefore preferable to use the request message  
20 only to request information that can subsequently also be displayed on the programming device.

Higher special purpose languages for sequencers (e.g., HIGRAPH, GRAPH 7, etc.) in turn use basic languages (e.g., AWL, KOP, FUP, etc.) to program the step enabling condition and the action to be performed in the corresponding step. For the user, it is important on the one hand to detect  
5 which step is currently active and, in addition, especially to look at the step enabling condition. For this purpose, it is generally sufficient to look at a few code addresses. The number of the relevant code addresses is relatively low compared to the size of the module. However, the relevant code addresses are usually scattered within the module and furthermore mixed in with the  
10 runtime system of the sequencer control. Based on the limited possibilities regarding data transfer between the automation device and the programming device, it would not make sense to monitor the entire module, since this might very quickly conflict with the upper limits of the data transfer capacity. In effect, of course, only the code addresses necessary to calculate the step  
15 enabling conditions are required. Consequently, only the actually required code addresses are transferred to the programming device in the present automation system.

A similar problem results in the programming of module calls with parameters. In the individual case, a plurality of *individual* instructions is  
20 required to transfer the parameter values. Consequently, only a few code addresses, which are distributed over a large area of the module, would actually be relevant for monitoring. In such a case, it is provided to accept the



values to be transferred to the module from the so-called stack. The filing of parameters on a stack when subroutines are called is one method for transferring parameters from a first program area to a second independent program area. It is apparent, therefore, that all values that are transferred into  
5 a module can be read from the stack, or even from a special instance data module, prior to the actual machine instruction to call the module. Analogously, all values that are transferred from the module can be read from the stack, or again from the instance data module, after the actual machine instruction to call the module. To monitor parameter transfer, two code  
10 addresses are required for each module call. In the individual case, however, several data addresses may be required per code address. The number of required data addresses depends on the number of input and output parameters.

Data flow control, also referred to as flow control in another context,  
15 is implemented in the following manner. After executing the modules associated with a cycle, the job initiated by the request message is blocked. This means that no further data is recorded, even if the automation device executes an instruction one additional time at a code address specified by the request message. At this time, the transmission of the requested data is  
20 triggered. This transmission takes place in parallel to the execution of the application program. Such a transmission process can extend over several cycles, if necessary, in which case the recorded data is transmitted in sections.

During transmission, the application program is only imperceptibly interfered with due to the sectional transmission in parallel to the execution of the application program and particularly, also due to the possible distribution of the data to be transmitted over several cycles. Only after the programming  
5 device has evaluated and has displayed the data thus received, does it send a message to release the previously blocked job, so that recording of the data specified by the request message can restart.

In the programming device, the application programs are translated into machine instructions of the automation device by means of specially  
10 provided software. These can be instructions that are directly executable by the processor of the automation device or an intermediate language that is first interpreted in the automation system or translated into instructions of the processor of the automation system by the firmware of the automation system. In principle, however, the emphasis in translation is on short and  
15 quickly executable code.

In an automation device using a commercially available processor or processors in the form of microcontrollers, the machine model of the automation device must be transferred to the memory model of the microcontroller. For monitoring, additional code must be inserted into the  
20 translated application program. This generally requires a special translation run. If translation takes place in the firmware of the automation device, translation and the sequencing-in of the resulting additional instructions

burdens the cycle time. If translation takes place in the programming device on the other hand, the cycle time is burdened by downloading and sequencing-in of the additional instructions. In either case, the present invention allows for the burden on the cycle time to be minimized by  
 5 appropriately selecting the code addresses to be monitored.

To prevent burdening of the cycle time when a monitoring job is blocked, the instructions required for monitoring can be masked. The requested data is thus recorded only if the corresponding instructions are not masked. After completion of the cycle, the job can be blocked by the  
 10 firmware unless this already occurs during recording. Finally, the recorded data is sent to the programming device.

In an automation device that interprets the machine instructions, the interpreter causes the machine model to be mapped to the real memory. For the following examination it is unimportant whether the interpreter is realized  
 15 by an application specific integrated circuit (ASIC) or by a program on a commercially available microcontroller. The firmware flags the instructions to be monitored, for which it requires at least 1 bit per instruction. When the interpreter meets a flagged instruction, it interrupts execution and records the requested data. Subsequently, execution is continued. As in translating  
 20 machines, this interruption must be maskable to prevent burdening the cycle time whenever the job is blocked. Analogous to translating machines, the job can be blocked by the firmware after completion of the cycle, unless this

already occurs during recording. Finally, the recorded data is sent to the programming device.

According to a further embodiment of the present invention, a monitoring job is provided with conditions to ensure that in modules, which are multiply executed in a cycle, a special execution can be monitored. Furthermore, it is possible to distinguish in a job what the response is to be in case of multiple execution of the instructions within a module. Such multiple execution occurs in the module if loops are created by jumps within the module. In the monitoring job, it may be specified, for example, whether or not the job is blocked by multiple execution. If it is not blocked in the second execution of a code address, the data for this code address is overwritten so that, as a result, the data of the last looping is available for actual transfer to the programming device. If, in contrast, a transfer were to occur with each looping, a high burden would be placed on the cycle time. The burden on the cycle time can be further reduced in that recording is blocked when multiple execution first occurs, so that only the data of the first looping is transferred.

The present invention also allows for a specification to contain several code addresses for each module and several data addresses for each code address. The code address defines an individual machine instruction. The data addresses can be used to address all memory areas present in the machine model, including the registers. This method minimizes the data to be recorded and transferred as well as the burden on the cycle time associated

with said recording and transfer. Said minimization is independent of the language in which the application program was created and the type of the target machine.

### **BRIEF DESCRIPTION OF THE DRAWINGS**

5           The invention is explained in more detail below with the aid of diagrammatic, exemplary embodiments in the drawing, in which:

FIG 1 is an illustration of a request message according to a preferred embodiment of the invention.

10           FIG 2 is an illustration of a receive message according to a preferred embodiment of the invention.

FIG 3 is an illustration of an embodiment of an automation system according to a preferred embodiment of the invention.

### **DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS**

15           According to FIG 1, the request message AT has at least a header segment AH, an entry data set segment AE, and one or several watch point segments AW1, AW2, ..., AWn. The header segment AH contains: (a) the job identification 1110; (b) a coding of module 1111 coded by type and number, to which the request message refers; (c) a trigger condition 1112; (d)  
20           an identifier regarding possible multiple executions 1113; and (e) the number of watch points without entry data set 1114.

The entry data set AE, which in contrast to the watch points shows the corresponding monitoring values, particularly the registers, prior to execution of the instruction, comprises: (a) a code address 1121; (b) register 1122; (c) an identifier with respect to the number of the subsequent data  
5 addresses 1123; and (d) the number of data addresses 1124 defined by the identifier 1123.

The structure of the watch point segments AW1, AW2, ..., AWn is analogous to that of the entry data set AE. Each of these watch point segments AW1, AW2, ..., AWn comprises a code address 1131-1, 1131-2,  
10 1131-n ("-1" references the first watch point segment AW1, "-2" correspondingly references the second watch point segment and "-n" the watch point segment AWn. In addition to the code address, every watch point segment contains: (a) data on the relevant register 1132; (b) an identifier with respect to the number of the subsequent data addresses 1133; and (c) the  
15 number of data addresses 1134 specified according to the identifier 1133.

In accordance with FIG 2, the receive message ET comprises at least the header EH, entry data EE, and watch point data EW1, EW2, ..., EWn. Specifically, the header EH comprises at least a job identifier 1210 and a job number 1211. The entry data comprises at least the code address 1221,  
20 register values 1222, and data values 1223. The watch point data EW1, EW2, ..., EWn analogously comprise at least the code address 1231, register values 1232, and data values 1233.

The job number 1211 is used to distinguish jobs that are present simultaneously. A unique job number 1211 is automatically assigned by the automation device when the job is allocated. The size of the segments of the receive message ET provided for the register values, 1222, 1232-1, 1232-2, ..., 1232-n depends on the type and number of the requested registers. Analogously, the size of the segments of the receive message ET provided for the data values 1233, 1233-1, 1233-2, ..., 1233-n depends on the type and amount of the requested data.

The receive message ET does not necessarily contain all code addresses of the request message AT, since possible jumps in the execution of the program to be monitored may have caused parts of said program not to be executed and certain code addresses specified in the request message AT not to be reached.

The entry data set EE is always present, but its code address 1221 does not necessarily coincide with the code address of the request 1121, since a jump over this first address may have occurred in the monitored cycle.

Figure 3 illustrates the automation system, and the flow of data traffic between the programming device 20 and the automation device 10, which are suitably linked for this purpose by means, for example, of a data transmission cable 40. One of ordinary skill in the art would also recognize that data transmission between the programming device 20 and the automation device 10 can also occur without a physical connection, i.e., transmission and signal

technology. As shown in Figure 3, sensors 30 provide data for monitoring of an external process to the automation device 10. The automation device 10 and programming device 20 can both comprise a monitoring device 60 for displaying process information and messages, etc., as well as a computer readable medium 50 for processing, and storage purposes.

As soon as the user starts monitoring, a request message AT is composed on the side of the programming device 20. This request message AT is transmitted to the automation device 10. The automation device 10 receives the request message AT, accepts it if the structure of the request message AT corresponds to the expected specifications, and assigns a job number 1211. The job is then flagged as blocked.

The automation device 10 sends a positive acknowledgment together with the job number 1211 to the programming device 20. When it receives the positive acknowledgment, the programming device marks the job as accepted.

A release message is then composed on the side of the programming device 20. This release message is transmitted to the automation device 10. On the side of the automation device 10, the job is marked as released upon receipt of the release message.

In response, the automation device 10 sends a positive acknowledgment to the programming device 20. On the side of the



programming device 20, the job is marked as released upon receipt of the positive acknowledgment.

If, during execution of the application program in the automation device 10, the trigger condition 1112 specified in the header AH of the request message AT is met at a given instant, the recording of the data, which the user selected for monitoring, begins. In the simplest case, the trigger condition consists of reaching one of the code addresses 1131-1, 1131-2, ..., 1131-n indicated in the request message AT, but more complex trigger conditions are feasible, e.g., "call path" or "open data blocks." If one module out of several other modules is used within a cycle, it is thus possible to monitor any special execution of the code addresses within this cycle. The call path describes a module hierarchy, which is necessary for the recording to begin by the automation device 10. The condition "open data blocks" requires that the indicated data blocks must be open at the start of recording. The conditions may be used alone or in combination.

If, in a simple trigger condition of the aforementioned type, for instance, the code address referenced by address 1131-1 of the reply message AT, is met during execution of the application program, the registers referenced by datum 1132-1, the number of data addresses referenced by datum 1133-1, and notably the data referenced by datum 1134-1 are recorded.

After said recording has been completed, a result message is constructed, which contains at least this recorded data. The corresponding job

is blocked and the result is then transmitted to the programming device 20 in the form of a result message ET.

On the side of the programming device 20, the result can now be displayed using the display device 60, or the results can be transmitted to an external device for viewing such as an external monitor or printer 70 for printing hard copies of the data. If the code address referenced by datum 1231-1 is displayed on the monitor 60, the register values 1232-1 and the data values 1233-1 can be displayed in connection with this code address.

After the results have been displayed on the monitor 60, or external device 70, a release message is again sent to the automation device 10 to release the job again. The automation device 10 transmits a positive acknowledgment to the programming device 20 as soon as the job is released. The programming device 20 marks the job as released when it receives said positive acknowledgment. Thus the conditions are restored in such a way that, when the trigger condition is reached next, recording starts again and the recorded data is finally transmitted to the programming device 20 for display. These steps are repeated until the user terminates the job.

The above description of the preferred embodiments has been given by way of example. From the disclosure given, those skilled in the art will not only understand the present invention and its attendant advantages, but will also find apparent various changes and modifications to the structures disclosed. It is sought, therefore, to cover all such changes and modifications

